

(19) World Intellectual Property Organization  
International Bureau(43) International Publication Date  
18 July 2002 (18.07.2002)

PCT

(10) International Publication Number  
WO 02/056559 A1(51) International Patent Classification<sup>7</sup>: H04L 27/34, I103M 13/00

(21) International Application Number: PCT/IB01/02556

(22) International Filing Date:  
14 December 2001 (14.12.2001)

(25) Filing Language: English

(26) Publication Language: English

(30) Priority Data:  
01200152.5 16 January 2001 (16.01.2001) EP

(71) Applicant: KONINKLIJKE PHILIPS ELECTRONICS N.V. [NL/NL]; Groenewoudseweg 1, NL-5621 BA Eindhoven (NL).

(72) Inventors: GOROKHOV, Alexei; Prof. Holstlaan 6, NL-5656 AA Eindhoven (NL). VAN DIJK, Marten, E.; Prof. Holstlaan 6, NL-5656 AA Eindhoven (NL). KOPPELAAR, Arie, G., C.; Prof. Holstlaan 6, NL-5656 AA Eindhoven (NL).

(74) Agent: DEGUELLE, Wilhelmus, H., G.; Internationaal Octrooibureau B.V., Prof. Holstlaan 6, NL-5656 AA Eindhoven (NL).

(81) Designated States (*national*): AE, AG, AL, AM, AT, AU, AZ, BA, BB, BG, BR, BY, BZ, CA, CH, CN, CO, CR, CU, CZ, DE, DK, DM, DZ, EC, EE, ES, FI, GB, GD, GE, GH, GM, HR, HU, ID, IL, IN, IS, JP, KE, KG, KP, KR, KZ, LC, LK, LR, LS, LT, LU, LV, MA, MD, MG, MK, MN, MW, MX, MZ, NO, NZ, PH, PL, PT, RO, RU, SD, SE, SG, SI, SK, SL, TJ, TM, TR, TT, TZ, UA, UG, UZ, VN, YU, ZA, ZW.(84) Designated States (*regional*): ARIPO patent (GH, GM, KE, LS, MW, MZ, SD, SL, SZ, TZ, UG, ZM, ZW), Eurasian patent (AM, AZ, BY, KG, KZ, MD, RU, TJ, TM), European patent (AT, BE, CH, CY, DE, DK, ES, FI, FR, GB, GR, IE, IT, LU, MC, NL, PT, SE, TR), OAPI patent (BF, BJ, CF, CG, CI, CM, GA, GN, GQ, GW, ML, MR, NE, SN, TD, TG).

## Published:

- with international search report
- entirely in electronic form (except for this front page) and available upon request from the International Bureau

For two-letter codes and other abbreviations, refer to the "Guidance Notes on Codes and Abbreviations" appearing at the beginning of each regular issue of the PCT Gazette.

(54) Title: BIT INTERLEAVED CODED MODULATION (BICM) MAPPING

0000	0101	1110	1011	0000	1011	1110	0101	0000	1001	1110	0111	0000	1010	0101	1001
1100	1001	0010	0111	1100	0111	0010	1001	1100	0101	0010	1011	1100	0110	1111	0011
1010	1111	0100	0001	1010	0001	0100	1111	0011	1111	1000	0100	1011	0001	1000	0100
0110	0011	1000	1101	0110	1101	1000	0011	1010	0110	0001	1101	0111	1101	0010	1110
A				B				C				D			

(57) Abstract: Described is a transmission system for transmitting a multilevel signal ( $X_k$ ) from a transmitter (10) to a receiver (20). The transmitter (10) comprises a mapper (16) for mapping an input signal ( $i_k$ ) according to a signal constellation onto the multilevel signal ( $X_k$ ). The receiver (20) comprises a demapper (22) for demapping the received multilevel signal ( $Y_k$ ) according to the signal constellation. The signal constellation comprises a number of signal points with corresponding labels. The signal constellation is constructed such that  $D_a > D_b$ , with  $D_b$  being the minimum of the Euclidean distances between all pairs of signal points whose corresponding labels differ in a single position, and with  $D_b$  being the minimum of the Euclidean distances between all pairs of signal points. By using this signal constellation a significantly lower error rate can be achieved than by using a prior-art signal constellation, in an iteratively decoded BICM (Bit Interleaved Coded Modulation) system.

WO 02/056559 A1

## BIT INTERLEAVED CODED MODULATION (BICM) MAPPING

The invention relates to a transmission system for transmitting a multilevel signal from a transmitter to a receiver.

The invention further relates to a transmitter for transmitting a multilevel signal, a receiver for receiving a multilevel signal, a mapper for mapping an interleaved  
 5 encoded signal according to a signal constellation onto a multilevel signal, a demapper for demapping a multilevel signal according to a signal constellation, a method of transmitting a multilevel signal from a transmitter to a receiver and to a multilevel signal.

In transmission systems employing so-called bit interleaved coded modulation (BICM) schemes a sequence of coded bits is interleaved prior to being encoding to channel  
 10 symbols. Thereafter, these channels symbols are transmitted. A schematic diagram of a transmitter 10 which may be used in such a transmission system is shown in Fig. 1. In this transmitter 10 a signal comprising a sequence of information bits  $\{b_k\}$  is encoded in a Forward Error Control (FEC) encoder 12. Next, the encoded signal  $\{c_k\}$  (i.e. the output of the  
 encoder 12) is supplied to an interleaver 14 which interleaves the encoded signal by  
 15 permuting the order of the incoming bits  $\{c_k\}$ . The output signal  $\{i_k\}$  of the interleaver 14 (i.e. the interleaved encoded signal) is then forwarded to a mapper 16 which groups the incoming bits into blocks of  $m$  bits and maps them to a symbol set consisting of  $2^m$  signal constellation points with corresponding labels. The resulting sequence of symbols  $\{x_k\}$  is a multilevel signal which is transmitted by the transmitter 10 over a memoryless fading  
 20 channel to a receiver 20 as shown in Fig. 2. In Fig. 1 the memoryless fading channel is modeled by the concatenation of a multiplier 17 and an adder 19. The memoryless fading channel is characterized by a sequence of gains  $\{\gamma_k\}$  which are applied to the transmitted multilevel signal by means of the multiplier 17. Furthermore, the samples of the transmitted multilevel signal are corrupted by a sequence  $\{n_k\}$  of Additive White Gaussian Noise  
 25 (AWGN) components which are added to the multilevel signal by means of the adder 19. This generic channel model fits, in particular, multicarrier transmission over a frequency selective channel, where the set of instances  $k = 1, \dots, N$  corresponds to  $N$  subcarriers. Therefore, it falls under the scope of the existing standards for broadband wireless (such as ETSI BRAN HIPERLAN/2, IEEE 802.11a and their advanced versions currently being in

standardization). The main distinguishing feature of BICM schemes is the interleaver 14 which spreads the adjacent encoded bits  $c_k$  over different symbols  $x_k$ , thereby providing the diversity of fading gains  $\gamma_k$  within a limited interval of the sequence  $\{c_k\}$  of coded bits. This yields a substantial improvement in the FEC performance in fading environments. (Pseudo-  
 5 )random interleavers may be used that for big block size  $N$  guarantee a uniform spreading and therefore a uniform diversity over the whole coded sequence. Alternatively, row-column interleavers may be used.

It is now assumed that the receiver 20 has a perfect knowledge of the fading gains  $\{\gamma_k\}$ . This assumption is valid as in practice these gains can be determined very  
 10 accurately (e.g. by means of pilot signals and/or training sequences). The standard decoding of a BICM-encoded signal has a mirror structure to the structure of the transmitter 10 as shown in Fig. 1. For each  $k$ , the received samples  $y_k$  and the fading gains  $\gamma_k$  are used to compute the so-called a posteriori probabilities (APP) of all  $2^m$  signal constellation points for  $x_k$ . These APP values are then demapped, i.e. transformed to reliability values of individual  
 15 bits of the  $k$ -th block. The reliability value of a bit may be computed as a log-ratio of the APP of this bit being 0 over the APP of this bit being 1, given the set of APP values of  $2^m$  constellation points for the  $k$ -th block. Sometimes the APP of a bit being 0 or 1 is replaced by the bitwise maximum likelihood (ML) metrics, i.e. the largest APP over the constellation points matching this bit value. In this way the numerical burden can be reduced. These  
 20 reliability values are deinterleaved and forwarded to a FEC decoder which estimates the sequence of information bits, e.g. by means of standard Viterbi decoding.

The main drawback of this standard decoding procedure, as compared to the (theoretically possible but impractical) optimal decoding comes from the fact that there is no  
 simultaneous use of the codeword structure (imposed by FEC) and the mapping structure.  
 25 Although the strictly optimal decoding is not feasible, the above observation gives rise to a better decoding procedure that is illustrated in the receiver 20 as shown in Fig. 2. The basic idea of this procedure is to iteratively exchange the reliability information between the demapper 22 and the FEC decoder 32. The iterative procedure starts with the standard demapping as described above. The reliability values  $\{L_k^{(\mu)}\}$  of the demapped bits, after  
 30 deinterleaving by a deinterleaver 26, serve as the inputs to a soft-input soft-output (SISO) decoder 32 which produces the (output) reliabilities  $\{L_k^{(c)}\}$  of the coded bits  $\{c_k\}$  that take into account the (input) reliabilities of the demapped bits and the FEC structure. The standard SISO decoders are maximum a posteriori (MAP) decoders, a simplified version of which is

known as a max-log-MAP (MLM) decoder. The difference between the inputs and the outputs of the SISO decoder 32 (often referred to as extrinsic information) is determined by a subtracter 30 and reflects the reliability increment which is the result of the code structure. This differential reliability is interleaved by an interleaver 28 and used as an a priori  
5 reliability during the next demapping iteration. In a similar way, the differential reliability is computed at the successive demapper output (by means of subtracter 24). This reliability represents a refinement due to the reuse of the mapping and signal constellation structure; it is used as an a priori reliability for the subsequent SISO decoding iteration. After the last iteration, the SISO output reliabilities  $\{L_k^{(b)}\}$  of the information bits are fed to a slicer 34 to  
10 produce final decisions  $\{\hat{b}_k\}$  on the information bits.

An important feature of the BICM scheme is the mapping of bits according to a signal constellation comprising a number of signal points with corresponding labels. The most commonly used signal constellations are PSK (BPSK, QPSK, up to 8-PSK) and 4-QAM, 16-QAM, 64-QAM and sometimes 256-QAM. Furthermore, the performance of the  
15 system depends substantially on the mapping design, that is, the association between the signal points of the signal constellation and their  $m$ -bit labels. The standard Gray mapping is optimal when the standard (non-iterative) decoding procedure is used. Gray mapping implies that the labels corresponding to the neighboring constellation points differ in the smallest possible number of  $m$  positions, ideally in only one. An example of a 16-QAM signal  
20 constellation with the Gray mapping ( $m = 4$ ) is shown in Fig. 3A. It can easily be seen that the labels of all neighboring signal points differ in exactly one position.

However, the use of alternative mapping designs or mappings may improve dramatically on the performance of BICM schemes whenever any version of the iterative decoding is exploited at the receiver. In European patent application number 0 948 140 an  
25 iterative decoding scheme as shown in Fig. 2 is used with what is referred to as anti-Gray encoding mapping. It is however not clear what is meant by this anti-Gray encoding mapping. In a paper entitled "Trellis-coded modulation with bit interleaving and iterative decoding" by X. Li and J. Ritcey, IEEE Journal on Selected Areas in Communications, volume 17, pages 715 to 724, April 1999, a noticeable performance improvement is achieved  
30 by means of a widely used mapping design known as the Set Partitioning (SP) mapping. An example of a 16-QAM signal constellation with the SP mapping is shown in Fig. 3B.

In European patent application number 0 998 045 and European patent application number 0 998 087 an information-theoretic approach to mapping optimization is disclosed. The core idea of this approach is to use a mapping that reaches the optimal value of the mutual information between the label bits and the received signal, averaged over the label bits. The optimal mutual information depends on the signal-to-noise ratio (SNR), the design number of iterations of the decoding procedure as well as on the channel model. The optimal value of the mutual information is the value that minimizes the resulting error rate. According to this approach, selection of the optimal mappings relies upon simulations of error rate performance versus the aforementioned mutual information for a given SNR, number of iterations and channel model, with the subsequent computation of mutual information for all candidate mappings. Such a design procedure is numerically intensive. Moreover, it does not guarantee optimal error rate performance of the system. Besides the standard Gray mapping, in these European patent applications two new mappings for 16-QAM signal constellations are proposed (which mappings will be referred to as optimal mutual information (OMI) mappings). 16-QAM signal constellations with these OMI mappings are shown in Figs. 3C and 3D.

It is an object of the invention to provide an improved transmission system for transmitting a multilevel signal from a transmitter to a receiver. This object is achieved in the transmission system according to the invention, said transmission system being arranged for transmitting a multilevel signal from a transmitter to a receiver, wherein the transmitter comprises a mapper for mapping an input signal according to a signal constellation onto the multilevel signal, and wherein the receiver comprises a demapper for demapping the received multilevel signal according to the signal constellation, wherein the signal constellation comprises a number of signal points with corresponding labels, and wherein  $D_a > D_f$ , with  $D_a$  being the minimum of the Euclidean distances between all pairs of signal points whose corresponding labels differ in a single position, and with  $D_f$  being the minimum of the Euclidean distances between all pairs of signal points. The Euclidean distance between two signal points is the actual ('physical') distance in the signal space between these two signal points. By using a signal constellation with a  $D_a$  which is larger than  $D_f$  a substantially lower error rate can be reached than by using any of the prior art signal constellations. Ideally,  $D_a$  is as large as possible (i.e.  $D_a$  has a substantially maximum value), in which case the error rate

is as low as possible.  $D_a$  is referred to as the effective free distance of the signal constellation and  $D_f$  is referred to as the exact free distance of the signal constellation.

It is observed that iterative decoding procedures approach the behavior of an optimal decoder when the SNR exceeds a certain threshold. This means that at a relatively high SNR (that ensures a good performance of the iterative decoding) one may assume that an optimal decoder is performing the decoding.

Consider an optimal decoder. In practice, trellis codes are used as FEC for noisy fading channels such as (concatenated) convolutional codes. A typical error pattern is characterized by a small number of erroneous coded bits  $\{c_k\}$  at error rates of potential interest. The number of erroneous coded bits is typically a small multiple of the free distance of the code; this number is only a small fraction of the total number of coded bits. The free distance of a code is the minimum number of bits (bit positions) in which two different codewords of the code can differ. Due to interleaving, these erroneous coded bits are likely to be assigned to different labels and therefore different symbols. More specifically, the probability of having only one erroneous coded bit per symbol approaches one along with the increase of the data block size.

Hence, the overall error rate (for error rates of potential interest) is improved when the error probability is decreased for such errors that at most one bit per symbol is corrupted. This situation can be reached by maximizing the minimum  $D_a$  of the Euclidean distances between all pairs of signal points whose corresponding labels differ in a single bit position.

In an embodiment of the transmission system according to the invention  $\overline{H}_1$  has a substantially minimum value, with  $\overline{H}_1$  being the average Hamming distance between all pairs of symbols corresponding to neighboring signal points. The Hamming distance between two labels is equal to the number of bits (bit positions) in which the labels differ. By this measure, an accurate decoding of the multilevel signal in the receiver is reached at a relatively small SNR. A typical feature of iterative decoding is a relatively poor performance up to some SNR threshold. After this threshold, the error rate of the iterative decoding approaches the performance of an optimal decoder quite soon, along with the increase of the SNR. It is therefore desirable to decrease this SNR threshold value. This threshold value depends on the starting point of the iterative procedure, i.e. on the distribution of the reliability values  $L_k^{(\mu)}$  provided by the demapper on the first iteration. The worst reliability values are due to the neighboring signal points, therefore the 'average' number of coded bits

that suffer from these poor reliabilities is proportional to the 'average' number of positions in which the labels, which correspond to the neighboring signal points, are different. In other words, the SNR threshold degrades (i.e. increases) along with the increase of the average Hamming distance between the labels that are assigned to the neighboring signal points.

- 5 Ideally,  $\overline{H_1}$  is as small as possible, i.e.  $\overline{H_1}$  has a minimum value, for which value of  $\overline{H_1}$  the SNR threshold will also be minimal.

The above object and features of the present invention will be more apparent from the following description of the preferred embodiments with reference to the drawings, wherein:

10

Fig. 1 shows a block diagram of a transmitter according to the invention,

Fig. 2 shows a block diagram of a receiver according to the invention,

Figs. 3A to 3D show prior-art 16-QAM signal constellations,

- 15 Fig. 4 shows graphs illustrating the packet error rate versus  $E_b/N_0$  (i.e. the SNR per information bit) for several 16-QAM mappings,

Fig. 5 shows graphs illustrating the bit error rate versus  $E_b/N_0$  for several 16-QAM mappings,

- 20 Fig. 6 shows graphs illustrating the packet error rate versus  $E_b/N_0$  for a standard 8-PSK signal constellation and for a modified 8-PSK signal constellation,

Fig. 7 shows graphs illustrating the bit error rate versus  $E_b/N_0$  for a standard 8-PSK signal constellation and for a modified 8-PSK signal constellation,

Figs. 8A to 8G show improved 16-QAM signal constellations,

Figs. 9A to 9C and Fig. 10 show improved 64-QAM signal constellations,

- 25 Figs. 11A and 11B show improved 256-QAM signal constellations,

Figs. 12A to 12C show improved 8-PSK signal constellations,

Fig. 13 shows a modified 8-PSK signal constellation.

- 30 In the Figs, identical parts are provided with the same reference numbers.

In Figs 8A to 8G, 9A to 9C, 10, 11A and 11B no horizontal I-axis and vertical Q-axis are shown. However, in these Figs. a horizontal I-axis and a vertical Q-axis must be considered to be present, which I-axis and Q-axis cross each other in the center of each Fig. (similar to the situation as shown in Figs. 3A to 3D).

The transmission system according to the invention comprises a transmitter 10 as shown in Fig. 1 and a receiver 20 as shown in Fig. 2. The transmission system may comprise further transmitters 10 and receivers 20. The transmitter 10 comprises a mapper 10 for mapping an input signal  $i_k$  according to a certain signal constellation onto a multilevel signal  $x_k$ . A multilevel signal comprises a number of groups of  $m$  bits which are mapped onto a real or complex signal space (e.g. the real axis or the complex plane) according to a signal constellation. The transmitter 10 transmits the multilevel signal  $x_k$  to the receiver 20 over a memoryless fading channel. The receiver 20 comprises a demapper 22 for demapping the received multilevel signal ( $y_k$ ) according to the signal constellation. The signal constellation comprises a number of signal points with corresponding labels. The (de)mapper is arranged for (de)mapping the labels to the signal constellation points such that  $D_a > D_f$ , with  $D_a$  being the minimum of the Euclidean distances between all pairs of signal points whose corresponding labels differ in a single position (these labels may be referred to as Hamming neighbors), and with  $D_f$  being the minimum of the Euclidean distances between all pairs of signal points. Such a mapping is referred to as a far neighbor (FAN) mapping.

Now the error rate performance of an iteratively decoded BICM scheme using the prior-art signal constellations as shown in Figs 3A to 3D will be compared with the error rate performance of an iteratively decoded BICM scheme using the 16-QAM FAN signal constellation as shown in Fig. 8E. The FEC coder 12 makes use of the standard 8-state rate (1/2) recursive systematic convolutional code with the feed-forward and feedback polynomials  $15_8$  and  $13_8$ , respectively. A sequence of 1000 information bits produces, after encoding, random interleaving and mapping, a set of  $N = 501$  symbols of 16-QAM that are transmitted over a Rayleigh channel with mutually independent gains  $\{\gamma_k\}$ . Note that this scenario fits a broadband multicarrier BICM scheme with a very selective multipath channel. At the receiver 20, an iterative decoding procedure is applied according to the scheme as shown in Fig. 2. In this example we use simplified (ML) reliability metrics for the demapping, along with a standard MLM SISO decoder. A pseudo-random uniform interleaver has been used. The simulation results are shown in Figs. 4 and 5. Fig. 4 shows the packet error rate (PER) versus  $E_b/N_o$  and Fig. 5 shows the bit error rate versus  $E_b/N_o$ . As expected, the signal constellation of Fig. 3A with the Gray mapping gives the worst results at desirably low error rates (see graphs 48 and 58). The state-of-the art signal constellation with SP mapping as shown in Fig. 3B improves substantially on this result (see graphs 44 and 54). The signal constellation with OMI mapping according to Fig. 3C (see graphs 46 and 56) has a poor packet error rate as compared to the signal constellation with SP mapping. However,



the signal constellation with OMI mapping of Fig. 3D (see graphs 42 and 52) improves substantially on the SP mapping. The signal constellation with FAN mapping as shown in Fig. 8E (see graphs 40 and 50) provides a 2 dB gain at low error rates (specifically at  $\text{PER} \leq 10^{-3}$ ) over the best of the prior-art signal constellations.

5           The effective free distance  $D_a$  is the minimum of the Euclidean distances taken over all pairs of signal points whose labels differ in one position only. Note that  $D_a$  is lower bounded by the exact free distance  $D_f$  which is the minimum Euclidean distance over all pairs of signal points.  $\overline{H}_1$  is defined as the average Hamming distance between the pairs of labels assigned to neighboring signal points (i.e. the signal points that are separated from each other  
10 by the minimum Euclidean distance  $D_f$ ). Now  $\overline{H}_l$  is defined as the average Hamming distance between all pairs of labels assigned to the  $l$ -th smallest Euclidean distance. By the  $l$ -th smallest Euclidean distance, the  $l$ -th element of an increasing sequence is meant, which sequence consists of all Euclidean distances between the signal points of a given constellation. Note that such a definition of  $\overline{H}_l$  is consistent with the definition of  $\overline{H}_1$ . In  
15 some cases, joint optimization of the first criterion (i.e. having a  $D_a$  which is as big as possible but at least larger than  $D_f$ ) and the second criterion (i.e. having a substantially minimum  $\overline{H}_1$ ) yields a set of solutions and some of them have different  $\overline{H}_l$  for some  $l > 1$ . In such cases, the set of solutions may be reduced in the following way. For each  $l$   
increasing from 1 to  $m$ , only the solutions that provide the minimum of  $\overline{H}_l$  are retained. This  
20 approach reduces the SNR threshold of the iterative decoding process.

All possible signal constellations may be grouped into classes of equivalent signal constellations. The signal constellations from the same equivalence class are characterized by the same sets of Euclidean and Hamming distances. Therefore all signal constellations of a given equivalence class are equally good for our purposes.

25           There are some obvious ways to produce an equivalent signal constellation to any given signal constellation. Moreover, the total number of equivalent signal constellations that may be so easily inferred from any given signal constellation is very big. The equivalence class of a given signal constellation is defined as a set of signal constellations that is obtained by means of an arbitrary combination of the following operations:

- 30           (a) choose an arbitrary binary  $m$ -tuple and add it (modulo 2) to all labels of the given signal constellation;
- (b) choose an arbitrary permutation of the positions of  $m$  bits and apply this permutation to all the labels;

(c) for any QAM constellation, rotate all signal points together with their labels by  $l\frac{\pi}{2}$ ,  $1 \leq l \leq 3$ ;

(d) for any QAM constellation, swap all signal points together with their labels upside down, or left to the right, or around the diagonals;

5 (e) for PSK, rotate all signal points together with their labels by an arbitrary angle.

A smart algorithm has been designed to accomplish the exhaustive classification of all possible signal constellations for 16-QAM for which  $D_a$  has a maximum value. This exhaustive search resulted in seven signal constellations which are shown in Figs. 8A to 8G. It is easy to show that all these signal constellations achieve the maximum possible effective free distance  $D_a$  which equals  $\sqrt{5}D_f$ . Note that all the prior-art signal constellations only achieve  $D_a = D_f$ .

In terms of the second criterion (i.e. having a substantially minimum  $\overline{H}_1$ ), the signal constellations of Figs. 8A to 8G have the respective  $\overline{H}_1$  values  
 15  $\left\{2\frac{1}{6}, 2\frac{1}{3}, 2\frac{1}{3}, 2\frac{1}{3}, 2\frac{1}{6}, 3, 2\frac{1}{3}\right\}$ . Note that the signal constellations of Figs. 8A and 8E yield the minimum of  $\overline{H}_1$ . Moreover, it can be shown that for 16-QAM,  $\overline{H}_1 = 2\frac{1}{6}$  is the minimum possible  $\overline{H}_1$  that may be achieved whenever  $D_a > D_f$ . Therefore, the signal constellations of Figs. 8A and 8E (and the signal constellations belonging to the equivalence classes thereof) jointly optimize both criteria under the condition  $D_a > D_f$ .

20 Since the total number of signal constellations grows very fast along with increasing  $m$  (For example, the total number of signal constellations is  $2.1 \cdot 10^{13}$ ,  $2.6 \cdot 10^{35}$  and  $1.3 \cdot 10^{89}$ , respectively, for  $m=4$ , 5 and 6, respectively) the exhaustive search for the best signal constellation is not feasible for  $m > 4$ . In such cases, an analytic construction should be found that allows to either simplify the exhaustive search or to restrict  
 25 ourselves to a limited set of signal constellations that contain 'rather good' ones.

As a matter of fact,  $2^m$ -QAM is almost the only signaling which is practically used for  $m \geq 4$ . For this signaling, with even  $m$ , we specify a family of linear signal constellations as follows:

Let  $m=2r$ , then the  $2^m$ -QAM signaling represents a regular two-dimensional grid with  $2^r$  points in the vertical and horizontal dimensions. A set of labels  $\{L_{i,j}\}_{1 \leq i,j \leq 2^r}$  is defined wherein  $L_{i,j}$  is a binary  $m$ -tuple that stands for the label of the signal point with the vertical coordinate  $i$  and the horizontal coordinate  $j$ . A signal constellation will be called  
 5 linear if and only if

$$L_{1,1} = O_m, \quad L_{i,j} = L_{1,1} \oplus L_{i,j}, \quad 1 \leq i, j \leq 2^r \quad (1)$$

wherein  $O_m$  is the all-zero  $m$ -tuple and  $\oplus$  denotes the modulo 2 addition.

This family of signal constellations is of interest because of the observation that all the signal constellations in the Figs. 8A-8G, except for the signal constellations in the  
 10 Figs. 8B and 8C, appear to be linear. These linear signal constellations as well as the following sub-family of linear signal constellations may also be constructed without applying the above mentioned (first and second) design criteria.

A sub-family of linear signal constellations can be obtained via the following equation:

$$15 \quad L_{i,1} = X_i A, \quad L_{j,1} = Y_j A, \quad 1 \leq i, j \leq 2^r \quad (2)$$

where  $\{X_i\}_{1 \leq i \leq 2^r}$  and  $\{Y_j\}_{1 \leq j \leq 2^r}$  are two arbitrary sets of binary  $m$ -tuples and  $A$  is an arbitrary  $m \times m$  matrix with binary inputs which is an invertible linear mapping in the  $m$ -dimensional linear space defined over the binary field with the modulo 2 addition.

The use of (2) allows to confine the exhaustive search over all possible linear  
 20 signal constellations to a search over the sets  $\{X_i\}, \{Y_j\}$ . For a given pair of sets  $\{X_i\}, \{Y_j\}$  and the desired  $D_a$ , a suitable  $A$  can easily be determined.

An exhaustive search within the sub-family (2) for 64-QAM led to the following results: 12 equivalence classes were found with  $D_a = \sqrt{20}D_f$ , which is the upper bound on  $D_a$  for 64-QAM. The further minimization of  $\bar{H}_1$  reduced this set to 3 equivalence  
 25 classes. All these classes achieve  $\bar{H}_1 = 2\frac{3}{14}$ . The corresponding signal constellations are shown in Figs. 9A to 9C.

Within the sub-family (2) signal constellations were searched that minimize  $\bar{H}_1$  under the condition  $D_a > D_f$ . For 64-QAM the theoretical minimum of  $\bar{H}_1$  is defined by the lower bound  $\bar{H}_1 \geq 2\frac{1}{14}$ . No signal constellations with  $\bar{H}_1 < 2\frac{3}{14}$  were found for

$D_a > \sqrt{17}D_f$ . For  $D_a = \sqrt{17}D_f$  there are 57 equivalence classes with  $\overline{H}_1 = 2\frac{1}{14}$ . Among those, a unique equivalence class was found that minimizes  $\overline{H}_2$ . This class achieves  $\overline{H}_2 = 2\frac{13}{49}$ ; it is shown in Fig. 10.

The following material on linear signal constellations is related to various  
 5 signal constellations for  $r > 3$ . For those cases, it was not possible to classify all possible  
 signal constellations nor to establish the upper bound on  $D_a$ . For 256-QAM, a limited search  
 within the sub-family (2) of linear signal constellations led us to a set of 16 equivalence  
 classes that achieve  $D_a = \sqrt{80}D_f$  and  $\overline{H}_1 = 2\frac{1}{10}$ . Among these 16 classes, we retained only  
 two classes that minimize  $\overline{H}_2$ , achieving thereby  $\overline{H}_2 = 2\frac{59}{75}$ . Their respective signal  
 10 constellations are given in Figs 11A and 11B.

---

For the general case of  $2^{2r}$ -QAM, a sub-set of (2) was designed with which  
 the effective free distance

$$D_a \geq \sqrt{5}2^{r-2}D_f \quad (3)$$

can be reached. This particular construction is described now. First of all, we restrict  
 15 ourselves to the sets of  $\{X_i\}_{1 \leq i \leq 2^r}$  and  $\{Y_j\}_{1 \leq j \leq 2^r}$  such that:

(a) the first  $r$  bits of  $X_i$  represent  $(i-1)$  in a binary notation whereas the  
 following  $r$  bits are zeros.

(b) the first  $r$  bits of  $Y_j$  are zeros whereas the following  $r$  bits represent  
 $(j-1)$  in a binary notation.

20 For sake of simplicity, this selection of  $\{X_i\}$  and  $\{Y_j\}$  will be referred to as  
 lexico-graphical. For 64-QAM ( $m = 6, r = 3$ ), the lexico-graphical sets are as follows:

$$\{X_1, X_2, \dots, X_8\} = \{000000, 001000, 010000, 011000, 100000, 101000, 110000, 111000\}$$

$$\{Y_1, Y_2, \dots, Y_8\} = \{000000, 000001, 000010, 000011, 000100, 000101, 000110, 000111\}$$

The advantage of the lexico-graphical selection is twofold. First, it ensures that  
 25  $(X_i + Y_j) \neq 0_m$  for all  $1 \leq i, j \leq 2^r$  except for  $i=j=1$ , thereby ensuring that all  $L_{i,j}$  are  
 different. Second, it allows to easily find  $A$  that satisfies (3). To do that, we need to ensure  
 that for every pair  $(L_{i,j}, L_{r,j})$  such that  $(L_{i,j} \oplus L_{r,j})$  has only one non-zero bit, the  
 corresponding signal points are at least  $D_a$  apart. Note that the total number of binary  $m$ -

tuples having only one non-zero bit is  $m$ . These labels are represented by the rows of the  $m \times m$  identity matrix  $I_m$  such that  $(I_m)_{ii} = 1$  for all  $i$  and the other elements of  $I_m$  are zeros. Due to the linearity conditions (1) and (2), we have  $(L_{i,j} \oplus L_{r,j'}) =$

$((X_i \oplus X_r) \oplus (Y_j \oplus Y_{j'}))A$  for all  $1 \leq i, j \leq 2^r$ . Once again due to linearity, the matrix  $A$  can

5 be uniquely defined by the equation

$$Z_m = Z A, \text{ where } Z = \{Z_1^T \dots Z_m^T\}^T \quad (4)$$

is a  $m \times m$  matrix with binary inputs which is an invertible linear mapping in the  $m$ -dimensional linear space defined over the binary field with the modulo 2 addition (here  $(\cdot)^T$  denotes matrix transpose). We need to select  $m$  linearly independent  $m$ -tuples (row vectors)

10  $Z_i$  such that (3) is satisfied.

Of interest are all possible signal constellations that satisfy (1), (2) and (4) with  $\{Z_i\}_{1 \leq i \leq m}$  chosen so as to meet (3). According to (4),  $A$  will be given by the inverse of  $Z$  which, along with the lexico-graphical selection of  $\{X_i, Y_j\}$ , (1) and (2), specifies a signal constellation for the desired equivalence class.

15 Let us specify one particular selection of  $\{Z_i\}_{1 \leq i \leq m}$  and show that (3) holds:

choose the set of  $\{Z_i\}$  as an arbitrary ordering of all possible  $m$ -tuples that have 2 or 3 non-zero entries of which 2 (mandatory) non-zero entries are always at the first and the  $(r+1)$ -st position. Check that there are exactly  $m$ -tuples and that all they are linearly independent so that  $Z$  is invertible. We now show that (3) holds. Assume that  $L_{i,j}, L_{r,j'}$  are two arbitrary labels

20 that differ in one position only. Consequently,

$$L_{i,j} \oplus L_{r,j'} = (I_m)_l \quad (5)$$

i.e. the  $l$ -th row of  $I_m$ , for some  $l$  from  $\{1, 2, \dots, m\}$ . According to (1) and (2), we may write

$$L_{i,j} \oplus L_{r,j'} = ((X_i \oplus X_r) \oplus (Y_j \oplus Y_{j'}))A \quad (6)$$

Taking into account (4), (5), (6) and the fact that  $A$  is invertible, we find

$$25 \quad (X_i \oplus X_r) \oplus (Y_j \oplus Y_{j'}) = Z_l \quad (7)$$

Recall that, according to the lexico-graphical ordering, all  $X_i(Y_j)$  have zeros within the first (last)  $r$  positions. Let us inspect the spacings between the pairs  $(i,j)$   $(i',j')$  of signal points that satisfy (7) with the aforementioned selection for  $\{Z_l\}$ . First, consider the single possible  $m$ -tuple with only 2 mandatory non-zero entries. Check that (7) yields

$$30 \quad (X_i \oplus X_r) = \overbrace{100\dots0}^{(r-1)\text{times}} \overbrace{00\dots0}^{(r)\text{times}}, \quad (Y_j \oplus Y_{j'}) = \overbrace{00\dots01}^{(r)\text{times}} \overbrace{00\dots0}^{(r-1)\text{times}}$$

$$(X_i \oplus X_{i'})$$

Note that the first  $r$ -bits of  $(X_i \oplus X_{i'})$  and the last  $r$ -bits of  $(Y_i \oplus Y_{i'})$  read  $2^{r-1}$  in binary notation. According to the lexico-graphical selection of  $\{X_i\}$  and  $\{Y_j\}$ , the corresponding signal points  $(i,j)$  ( $i',j'$ ) have vertical and horizontal offsets of  $2^{r-1}$  positions. The resulting Euclidean distance between these points is composed of vertical and horizontal distances of

5  $2^{r-1} D_f$ .

Now, consider all  $m$ -tuples with 3 non-zero entries such that the third (non-mandatory) entry is one of the first  $r$  entries. We have

$$(X_i \oplus X_{i'}) = 1 \overbrace{??...?}^{(r-1)\text{times}} \overbrace{00...0}^{(r)\text{times}}, \quad (Y_j \oplus Y_{j'}) = \overbrace{00...01}^{(r)\text{times}} \overbrace{00...0}^{(r-1)\text{times}},$$

where  $1??...?$  has one non-zero entry within the last  $(r-1)$  entries. Using again the properties of the lexico-graphical selection, one can show that this yields a vertical offset of at least

10  $\frac{1}{2} 2^{r-1} = 2^{r-2}$  between the signal points  $(i,j)$  ( $i',j'$ ) whereas the horizontal offset remains  $2^{r-1}$ .

Clearly, the role of vertical/horizontal offsets exchange when we consider such  $Z_i$  that the third (non-mandatory) entry is one of the last  $r$  entries.

We see that, in all situations, the Euclidean distance between the signal points, whose labels differ in one position only, is composed of vertical and horizontal distances so

15 that one of them equals  $2^{r-1} D_f$  and the other one is not less than  $2^{r-2} D_f$ . Consequently, the minimum of the total Euclidean distance between such points satisfies

$$D_a \geq \sqrt{(2^{r-1} D_f)^2 + (2^{r-2} D_f)^2} = 2^{r-2} D_f \sqrt{2^2 + 1} = \sqrt{5} 2^{r-2} D_f \quad (8)$$

The non-linear family of signal constellation classes described hereafter may be seen as an extension of the linear family (1). This family comes from the equivalence classes (b) and (c) of all possible optimal classes for 16-QAM, (see Figs. 8A to 8G) that do not fall within the linear family. We noticed that the equivalence classes (b) and (c) may be regarded as being part of the family defined below.

20

Let  $S$  be a set of binary  $m$ -tuples that is closed under the (modulo 2) addition. We define an extension of the family of Fig. 8 as a collection of all equivalence classes of signal constellations having a set of labels  $\{L_{i,j}\}_{1 \leq i,j \leq 2^r}$  satisfying

25

$$L_{i,j} = 0_m, \quad L_{i,j} = L_{i,1} \oplus L_{1,j} \oplus f(L_{i,1} \oplus L_{1,j}), \quad 1 \leq i,j \leq 2^r \quad (9)$$

where  $f$  is a mapping from the set of  $m$ -tuples into itself such that firstly, for any  $m$ -tuple  $x$  from  $S$ ,  $f(x)$  is also in  $S$  and secondly,  $f(x) = f(y)$  for any  $m$ -tuples  $x, y$  such that  $(x \oplus y)$  is in  $S$ .

For 8-PSK, an exhaustive search was used to find the set of appropriate signal constellations. Apparently, there exist only three equivalence classes that satisfy  $D_a > D_f$ . These classes achieve  $D_a \approx 1.84776D_f$  and one of those has  $\overline{H}_1 = 2\frac{1}{2}$  while the remaining two achieve  $\overline{H}_1 = 2\frac{1}{4}$ . The corresponding signal constellations are shown in Figs 12A to 12C.

The success of the new strategy is based on the fact that coded bits are interleaved in such a way that the erroneous bits stemming from (typical) error events end up in different labels with a high probability. This property is ensured statistically when a random interleaver is used with a very big block size  $N$ . However, the probability of having more than one erroneous bit per label/symbol is different from zero when  $N$  is finite.

This observation leads to the following undesirable effect: the error floor (*i.e.*, the error rate flattening region) will be limited by a non-negligible fraction of error events that are characterized by more than one erroneous bit per label. In such cases, the potential gain due to high  $D_a$  will not be realized.

There exists a simple way to overcome the impact of such undesirable error events: the interleaver 14 should ensure that, for every label, the smallest number of the trellis sections (of the underlying FEC) between all pairs of the channel bits contributing to this label, is not less than a certain  $\delta > 0$ .

Such a design criterion ensures that a single error event may result in multiple erroneous bits per label if and only if this error event spans at least  $\delta$  trellis sections. For big  $\delta$ , the corresponding number of erroneous bits of this error event approximately equals  $(\delta/2R)$ , where  $R$  is the FEC rate. By choosing  $\delta$  big enough, we increase the Hamming distance of such undesirable error events thereby making them virtually improbable. Thus, choosing  $\delta$  big allows us to control the error floor irrespectively to the block size  $N$ . In our simulations a uniform random interleaver was used that satisfies this design criterion with  $\delta \geq 25$ .

The following result is based on our earlier observation that the effective free distance  $D_a$  of a BICM scheme may be substantially bigger than the exact free distance  $D_f$ ,

provided that FEC with interleaving and an appropriate signal constellation (i.e. having a  $D_a$  which is larger than  $D_f$ ) is used. Hence, it makes sense to design signal constellations that aim at increasing  $D_a$  rather than  $D_f$ .

This is supported by the following example. A new signal constellation is  
 5 derived from the standard 8-PSK signal constellation. Let us consider an instance of the new strategy represented by the signal constellation as depicted in Fig. 12C. A standard 8-PSK signal constellation is characterized by  $D_a^{8-PSK} = (1 - \cos(\pi/4))^{-1/2} D_f^{8-PSK} \approx 1.84776 D_f^{8-PSK}$ . It can easily be seen that this minimum distance is defined by the distances between the signal points within the pairs labeled (000,001), (110,111), (100,101) and (010,011). Indeed, these  
 10 are the only pairs such that the signal points are separated by the (minimum) rotation of  $(\pi/2)$  and their labels differ in one position. Note that  $D_a$  may be increased, e.g., by simply rotating (note that rotation preserves a highly desirable constant envelope property of PSK) the signal points labeled {001,111,101,011} (i.e. the second label within each pair) leftwise with a rotation angle  $\theta$ . An improved signal constellation with  $\theta = (3\pi/32)$  is shown in Fig.  
 15 13, wherein the empty bullets stand for the original places of the rotated points. This signal constellation achieves

$$D_a = \sqrt{((1 + \sin(3\pi/32))/(1 - \cos(\pi/4)))} D_f^{8-PSK} = \sqrt{(1 + \sin(3\pi/32))} D_a^{8-PSK} \approx 1.135907 D_a^{8-PSK}$$

In Figs. 6 and 7 the performance of the modified 8-PSK signal constellation of Fig. 13 (see graphs 60 and 70) is compared with the standard 8-PSK signal constellation of  
 20 Fig. 12C (see graphs 62 and 72). Note that the modified 8-PSK signal constellation leads to a slight degradation of at most 0.2dB at low SNR. This degradation is compensated by a gain of around 1dB at higher SNR. The modified 8-PSK signal constellation shows better performance at packet error rates below  $10^{-2}$ .

The scope of the invention is not limited to the embodiments explicitly  
 25 disclosed. The invention is embodied in each new characteristic and each combination of characteristics. Any reference signs do not limit the scope of the claims. The word "comprising" does not exclude the presence of other elements or steps than those listed in a claim. Use of the word "a" or "an" preceding an element does not exclude the presence of a plurality of such elements.



## CLAIMS:

1. A transmission system for transmitting a multilevel signal ( $x_k$ ) from a transmitter (10) to a receiver (20), the transmitter (10) comprising a mapper (16) for mapping an input signal ( $i_k$ ) according to a signal constellation onto the multilevel signal ( $x_k$ ), the receiver (20) comprising a demapper (22) for demapping the received multilevel signal ( $y_k$ ) according to the signal constellation, wherein the signal constellation comprises a number of signal points with corresponding labels, and wherein  $D_a > D_f$ , with  $D_a$  being the minimum of the Euclidean distances between all pairs of signal points whose corresponding labels differ in a single position, and with  $D_f$  being the minimum of the Euclidean distances between all pairs of signal points.
2. The transmission system according to claim 1, wherein  $D_a$  has a substantially maximum value.
3. The transmission system according to claim 1 or 2, wherein  $\overline{H_1}$  has a substantially minimum value, with  $\overline{H_1}$  being the average Hamming distance between all pairs of labels corresponding to neighboring signal points.
4. The transmission system according to claim 1 or 2, wherein the signal constellation is a 16-QAM signal constellation as depicted in any one of the Figs. 8A to 8G or an equivalent signal constellation thereof.
5. The transmission system according to claim 1 or 2, wherein the signal constellation is a 64-QAM signal constellation as depicted in any one of the Figs. 9A to 9C and 10 or an equivalent signal constellation thereof.
6. The transmission system according to claim 1 or 2, wherein the signal constellation is a 256-QAM signal constellation as depicted in any one of the Figs. 11A and 11B or an equivalent signal constellation thereof.

7. The transmission system according to claim 1 or 2, wherein the signal constellation is a 8-PSK signal constellation as depicted in any one of the Figs. 12A to 12C or an equivalent signal constellation thereof.
- 5 8. A transmitter (10) for transmitting a multilevel signal ( $x_k$ ), the transmitter (10) comprising a mapper (16) for mapping an input signal ( $i_k$ ) according to a signal constellation onto the multilevel signal ( $x_k$ ), wherein the signal constellation comprises a number of signal points with corresponding labels, and wherein  $D_a > D_f$ , with  $D_a$  being the minimum of the Euclidean distances between all pairs of signal points whose corresponding labels differ in a  
10 single position, and with  $D_f$  being the minimum of the Euclidean distances between all pairs of signal points.
9. The transmitter (10) according to claim 8, wherein  $D_a$  has a substantially  
15 maximum value.
10. A transmitter (10) according to claim 8 or 9, wherein  $\overline{H_1}$  has a substantially minimum value, with  $\overline{H_1}$  being the average Hamming distance between all pairs of labels corresponding to neighboring signal points.
- 20 11. A receiver (20) for receiving a multilevel signal ( $y_k$ ), the receiver (20) comprising a demapper (22) for demapping the multilevel signal ( $y_k$ ) according to a signal constellation, wherein the signal constellation comprises a number of signal points with corresponding labels, and wherein  $D_a > D_f$ , with  $D_a$  being the minimum of the Euclidean distances between all pairs of signal points whose corresponding labels differ in a single  
25 position, and with  $D_f$  being the minimum of the Euclidean distances between all pairs of signal points.
12. The receiver (20) according to claim 11, wherein  $D_a$  has a substantially  
30 maximum value.
13. The receiver (20) according to claim 11 or 12, wherein  $\overline{H_1}$  has a substantially minimum value, with  $\overline{H_1}$  being the average Hamming distance between all pairs of labels corresponding to neighboring signal points.

14. A mapper (16) for mapping an input signal ( $i_k$ ) according to a signal constellation onto a multilevel signal ( $x_k$ ), wherein the signal constellation comprises a number of signal points with corresponding labels, and wherein  $D_a > D_f$ , with  $D_a$  being the minimum of the Euclidean distances between all pairs of signal points whose corresponding labels differ in a single position, and with  $D_f$  being the minimum of the Euclidean distances between all pairs of signal points.
15. The mapper (16) according to claim 14, wherein  $D_a$  has a substantially maximum value.
16. The mapper (16) according to claim 14 or 15, wherein  $\overline{H_1}$  has a substantially minimum value, with  $\overline{H_1}$  being the average Hamming distance between all pairs of labels corresponding to neighboring signal points.
17. A demapper (22) for demapping a multilevel signal ( $y_k$ ) according to a signal constellation, wherein the signal constellation comprises a number of signal points with corresponding labels, and wherein  $D_a > D_f$ , with  $D_a$  being the minimum of the Euclidean distances between all pairs of signal points whose corresponding labels differ in a single position, and with  $D_f$  being the minimum of the Euclidean distances between all pairs of signal points.
18. The demapper (22) according to claim 17, wherein  $D_a$  has a substantially maximum value.
19. The demapper (22) according to claim 17 or 18, wherein  $\overline{H_1}$  has a substantially minimum value, with  $\overline{H_1}$  being the average Hamming distance between all pairs of labels corresponding to neighboring signal points.
20. A method of transmitting a multilevel signal ( $x_k$ ) from a transmitter (10) to a receiver (20), the method comprising the steps of :  
- mapping an input signal ( $i_k$ ) according to a signal constellation onto the multilevel signal ( $x_k$ ),

- transmitting the multilevel signal ( $x_k$ ),
  - receiving the multilevel signal ( $y_k$ ) and
  - demapping the multilevel signal ( $y_k$ ) according to the signal constellation, wherein the signal constellation comprises a number of signal points with corresponding labels, and
- 5 wherein  $D_a > D_f$ , with  $D_a$  being the minimum of the Euclidean distances between all pairs of signal points whose corresponding labels differ in a single position, and with  $D_f$  being the minimum of the Euclidean distances between all pairs of signal points.

21. The method according to claim 20, wherein  $D_a$  has a substantially maximum  
10 value.

22. The method according to claim 20 or 21, wherein  $\overline{H_1}$  has a substantially minimum value, with  $\overline{H_1}$  being the average Hamming distance between all pairs of labels corresponding to neighboring signal points.

15 23. A multilevel signal, the multilevel signal being the result of a mapping of an input signal ( $i_k$ ) according to a signal constellation, wherein the signal constellation comprises a number of signal points with corresponding labels, and wherein  $D_a > D_f$ , with  $D_a$  being the minimum of the Euclidean distances between all pairs of signal points whose  
20 corresponding labels differ in a single position, and with  $D_f$  being the minimum of the Euclidean distances between all pairs of signal points.

24. The multilevel signal according to claim 23, wherein  $D_a$  has a substantially maximum value.

25

25. The multilevel signal according to claim 23 or 24, wherein  $\overline{H_1}$  has a substantially minimum value, with  $\overline{H_1}$  being the average Hamming distance between all pairs of labels corresponding to neighboring signal points.

30 26. The multilevel signal according to claim 23 or 24, wherein the signal constellation is a 16-QAM signal constellation as depicted in any one of the Figs. 8A to 8G or an equivalent signal constellation thereof.

27. The multilevel signal according to claim 23 or 24, wherein the signal constellation is a 64-QAM signal constellation as depicted in any one of the Figs. 9A to 9C and 10 or an equivalent signal constellation thereof.
- 5 28. The multilevel signal according to claim 23 or 24, wherein the signal constellation is a 256-QAM signal constellation as depicted in any one of the Figs. 11A and 11B or an equivalent signal constellation thereof.
29. The multilevel signal according to claim 23 or 24, wherein the signal  
10 constellation is a 8-PSK signal constellation as depicted in any one of the Figs. 12A to 12C or an equivalent signal constellation thereof.

1/14

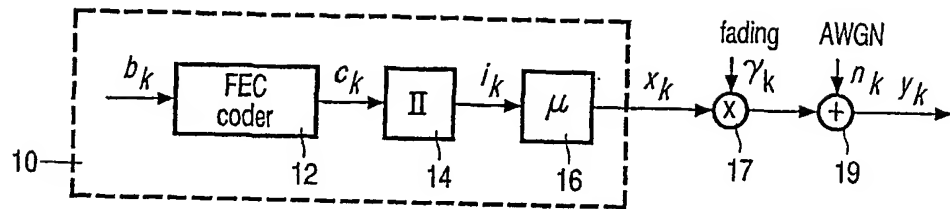


FIG. 1

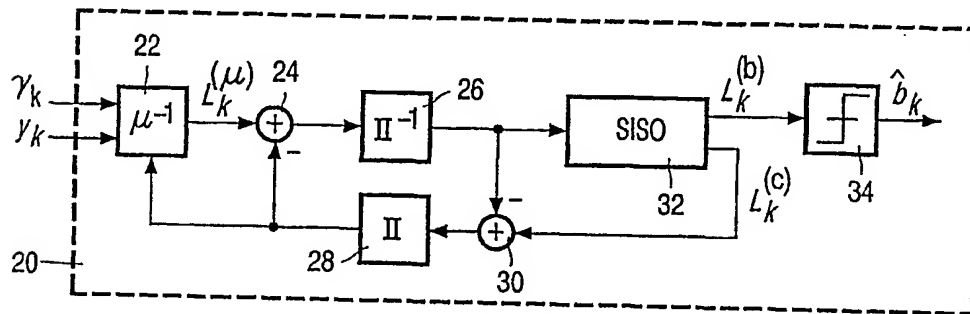


FIG. 2

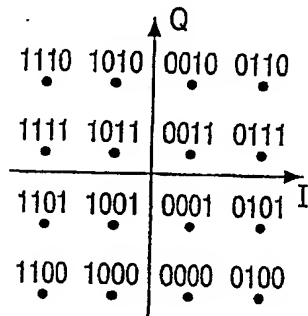


FIG. 3A

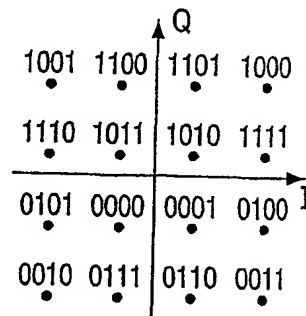


FIG. 3B

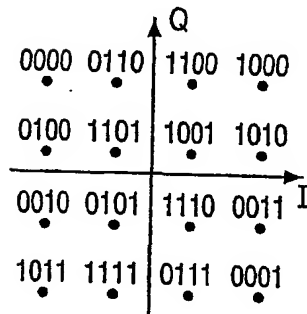


FIG. 3C

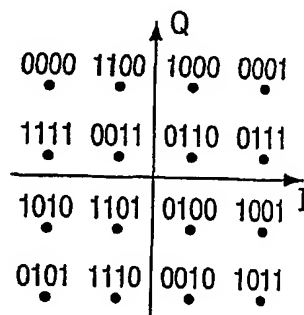


FIG. 3D

2/14

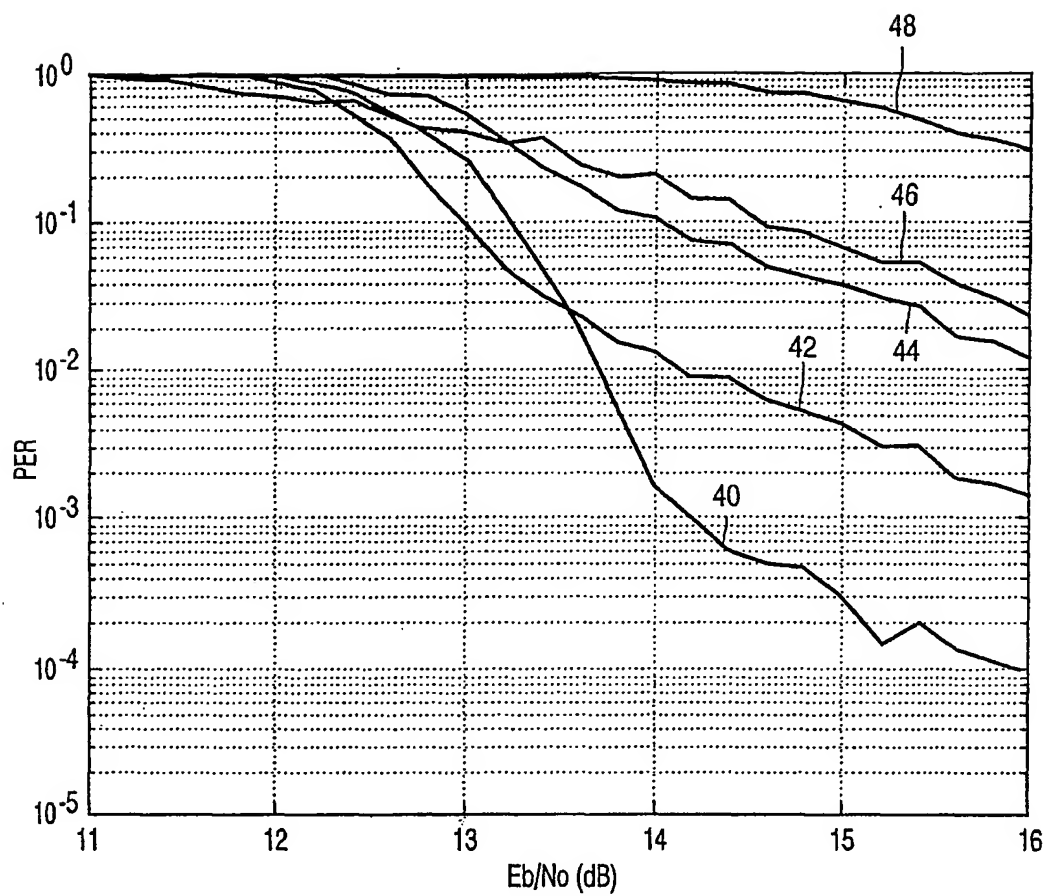


FIG. 4

3/14

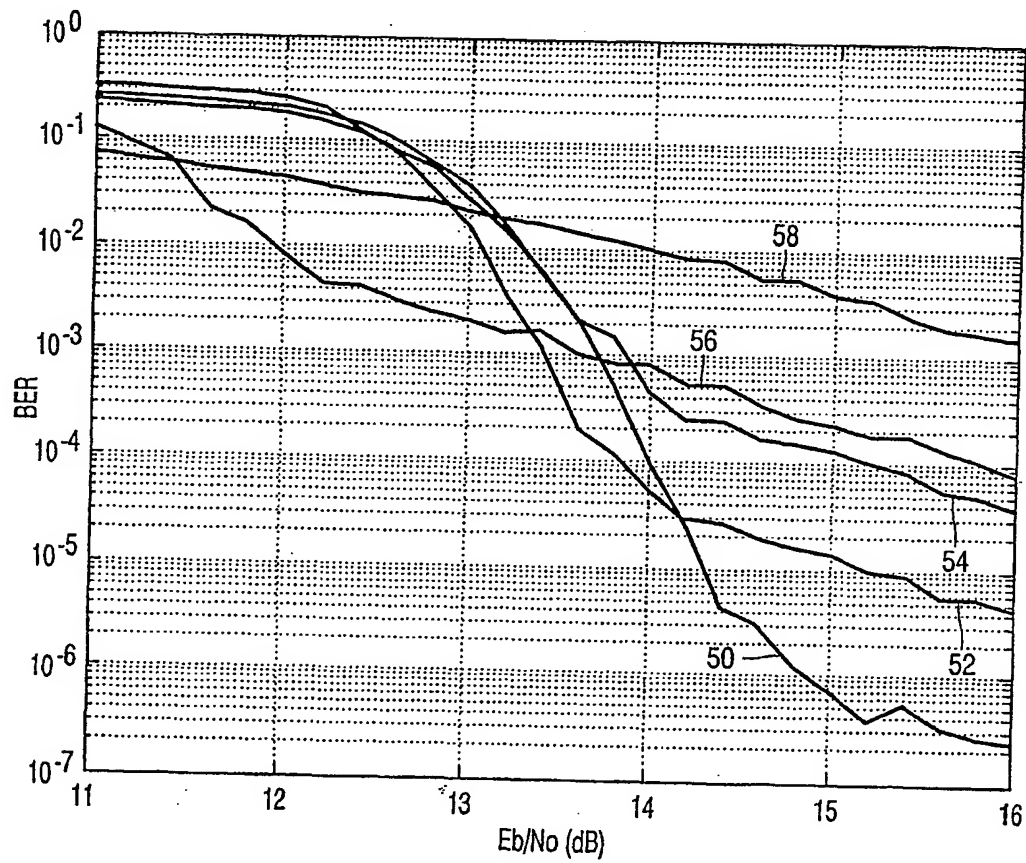


FIG. 5



4/14

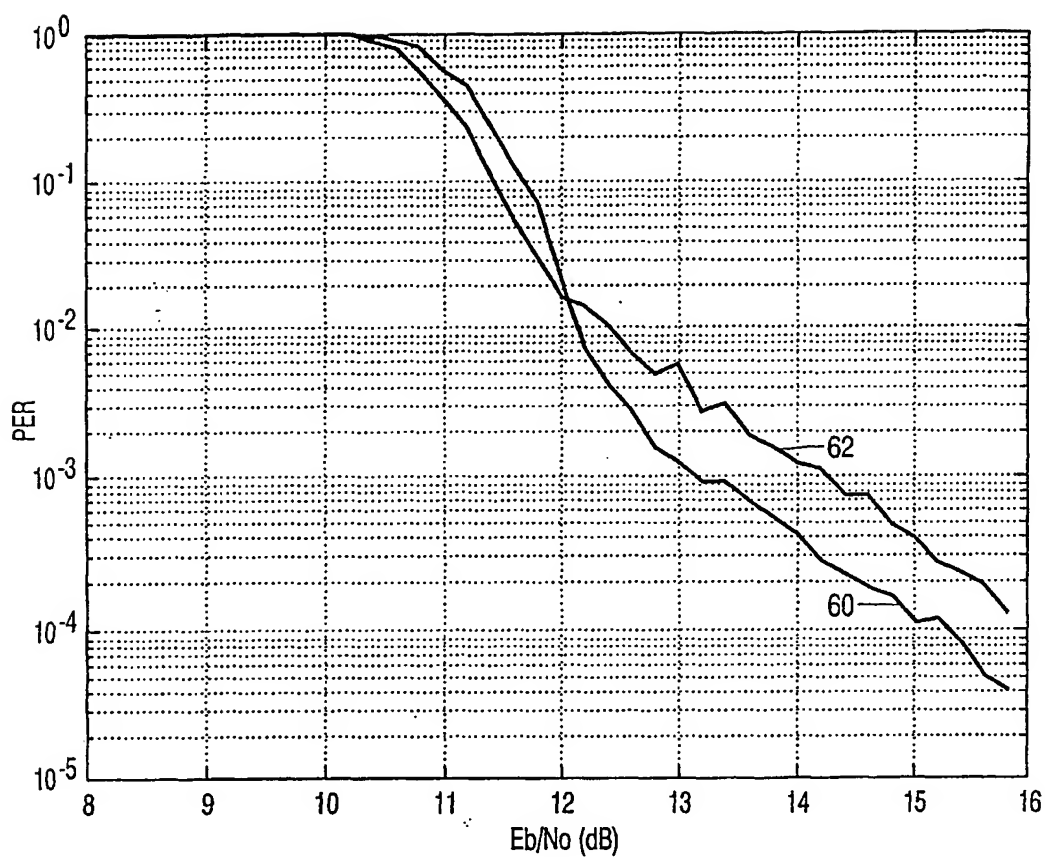


FIG. 6

5/14

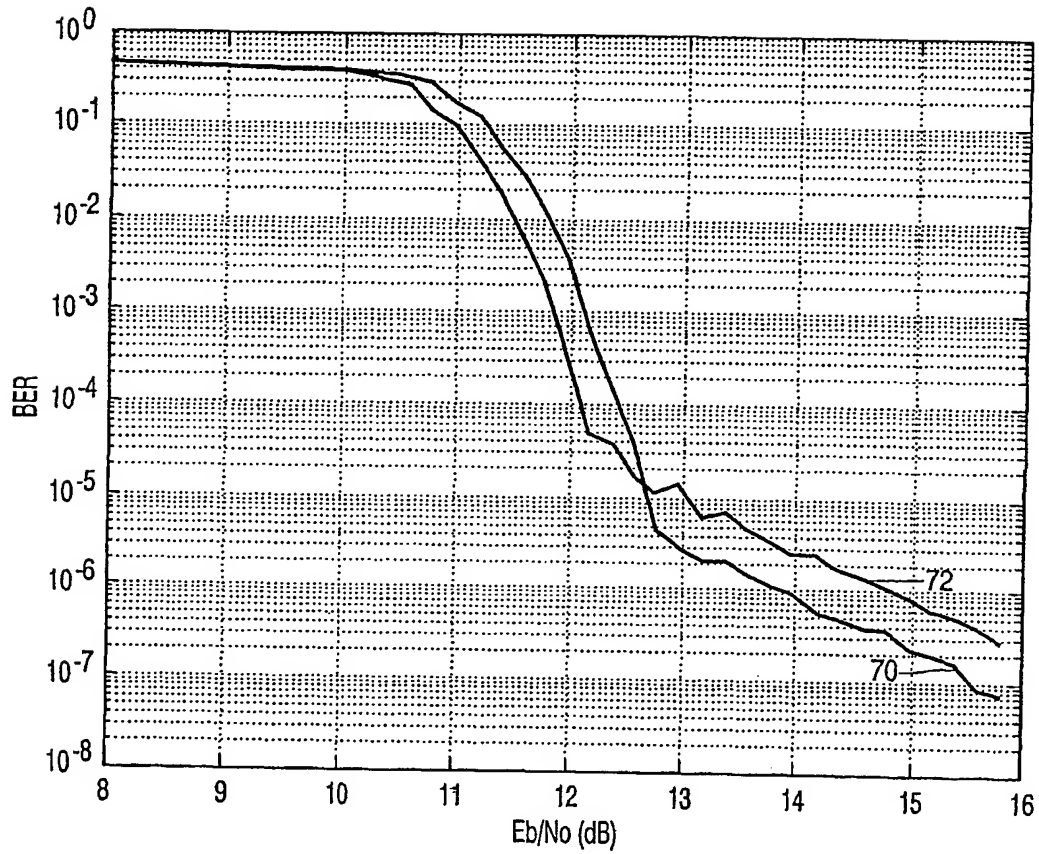


FIG. 7

6/14

0000	0101	1110	1011	0000	1011	1110	0101	0000	1010	0101	1001
1100	1001	0010	0111	1100	0111	0010	1001	1100	0110	1111	0011
1010	1111	0100	0001	1010	0001	0100	1111	0011	0001	1000	0100
0110	0011	1000	1101	0110	1101	1000	0011	1010	1101	0010	1110

FIG. 8A

FIG. 8B

FIG. 8C

FIG. 8D

7/14

0000	1010	0011	1001	0000	1101	0110	1011	0000	1001	1100	0101
1100	0110	1111	0101	1110	0011	1000	0101	1110	0111	0010	1011
1011	0001	1000	0010	1001	0100	1111	0010	1101	0100	0001	1000
0111	1101	0100	1110	0111	1010	0001	1100	0011	1010	1111	0110

FIG. 8E

FIG. 8F

FIG. 8G

8/14

000000 •	011000 •	010100 •	001100 •	100100 •	111100 •	110000 •	101000 •
010010 •	001010 •	000110 •	011110 •	110110 •	101110 •	100010 •	111010 •
010001 •	001001 •	000101 •	011101 •	110101 •	101101 •	100001 •	111001 •
000011 •	011011 •	010111 •	001111 •	100111 •	111111 •	110011 •	101011 •
110100 •	101100 •	100000 •	111000 •	010000 •	001000 •	000100 •	011100 •
100110 •	111110 •	110010 •	101010 •	000010 •	011010 •	010110 •	001110 •
100101 •	111101 •	110001 •	101001 •	000001 •	011001 •	010101 •	001101 •
110111 •	101111 •	100011 •	111011 •	010011 •	001011 •	000111 •	011111 •

FIG. 9A

9/14

000000 •	011000 •	010100 •	001100 •	110001 •	101001 •	100101 •	111101 •
010010 •	001010 •	000110 •	011110 •	100011 •	111011 •	110111 •	101111 •
010001 •	001001 •	000101 •	011101 •	100000 •	111000 •	110100 •	101100 •
000011 •	011011 •	010111 •	001111 •	110010 •	101010 •	100110 •	111110 •
100001 •	111001 •	110101 •	101101 •	010000 •	001000 •	000100 •	011100 •
110011 •	101011 •	100111 •	111111 •	000010 •	011010 •	010110 •	001110 •
110000 •	101000 •	100100 •	111100 •	000001 •	011001 •	010101 •	001101 •
100010 •	111010 •	110110 •	101110 •	010011 •	001011 •	000111 •	011111 •

FIG. 9B

10/14

000000 •	011000 •	010100 •	001100 •	110010 •	101010 •	100110 •	111110 •
000011 •	011011 •	010111 •	001111 •	110001 •	101001 •	100101 •	111101 •
010010 •	001010 •	000110 •	011110 •	100000 •	111000 •	110100 •	101100 •
110011 •	101011 •	100111 •	111111 •	000001 •	011001 •	010101 •	001101 •
100010 •	111010 •	110110 •	101110 •	010000 •	001000 •	000100 •	011100 •
100001 •	111001 •	110101 •	101101 •	010011 •	001011 •	000111 •	011111 •
110000 •	101000 •	100100 •	111100 •	000010 •	011010 •	010110 •	001110 •
010001 •	001001 •	000101 •	011101 •	100011 •	111011 •	110111 •	101111 •

FIG. 9C

11/14

000000 •	011000 •	001100 •	010100 •	110010 •	101010 •	111110 •	100110 •
010010 •	001010 •	011110 •	000110 •	100000 •	111000 •	101100 •	110100 •
000011 •	011011 •	001111 •	010111 •	110001 •	101001 •	111101 •	100101 •
110011 •	101011 •	111111 •	100111 •	000001 •	011001 •	001101 •	010101 •
100010 •	111010 •	101110 •	110110 •	010000 •	001000 •	011100 •	000100 •
110000 •	101000 •	111100 •	100100 •	000010 •	011010 •	001110 •	010110 •
100001 •	111001 •	101101 •	110101 •	010011 •	001011 •	011111 •	000111 •
010001 •	001001 •	011101 •	000101 •	100011 •	111011 •	101111 •	110111 •

FIG. 10



12/14

```

00000000 01100000 00110000 01010000 01001000 00101000 01111000 00011000 00011000 11110001 10010001 10010001 10111001 10111001 11011001
01000100 00100100 01110100 00010100 00001100 01101100 00111100 01011100 01011100 11100101 10101010 11001101 10101101 11111101 10011101
00000110 01100110 01010110 01001110 01001110 00101110 01111110 00011110 11000111 10100111 11110111 10010111 10001111 11011111 10111111
01000010 01110010 00010010 00001010 01101010 00111010 00111010 01011010 01011010 11100011 10110011 11010011 11001011 10101011 10011011
01000001 01110001 00010001 00001001 01101001 00111001 00111001 01011001 01011001 11100000 10110000 11010000 10010000 11111000 10011000
00000101 01100101 01010101 01001101 01001101 01111101 00011101 00011101 11000100 10100100 11110100 10010100 10001100 11101100 10111100
01000111 00100111 00010111 00001111 01101111 00111111 00111111 01011111 10000110 11100110 10101110 11001110 10101110 11111110 10011110
00000011 01100011 01010011 01001011 00101011 01111011 00011011 11000010 10100010 11110010 10010010 10001010 11101010 10111010 10111010
10000001 11100001 10110001 11010001 11001001 10101001 11111001 10011001 01000000 00100000 01110000 00010000 01101000 00111000 01011000
11000101 10100101 11110101 10010101 10001101 11101101 10111101 10111101 00000100 01110010 00110100 01010100 01001100 01111100 00011100
10000111 11100111 10110111 11010111 11001111 10101111 11111111 10011111 01000110 00100110 00010110 00001110 01101110 00111110 01011110
11000011 10100011 11110011 10010011 10001011 11101011 10111011 10111011 00000010 011100010 00110010 01010010 01001010 01111010 00011010
11000000 10100000 11110000 10010000 10001000 11101000 10111000 10111000 00000001 01110000 00010001 01010001 01001001 01111001 00011001
10000100 11110010 10110100 11010100 11001100 10101100 11111100 10011100 01000101 00100101 01110101 00010101 00001101 01101101 01011101
11000110 10100110 11110110 10010110 10001110 11101110 10111110 10111110 00000111 01100111 00110111 01010111 01001111 01111111 00011111
10000010 11100010 10110010 11010010 11001010 11001010 10101010 11111010 01000011 00100011 01110011 00010011 00001011 01101011 01011011

```

FIG. 11A

13/14

```

00000000 01100000 00110000 01010000 01001000 00101000 01111000 00011000 10001000 11011000 10100000 10100000 10010000
01000100 00100100 01110100 00010100 00001100 01101100 00111100 01011100 11001100 10101100 11111100 10001100 11100100
00000110 01100110 00110110 01010110 01001110 00101110 01111110 00011110 10001110 11101110 10111110 10100110 10010110
01000010 00100010 01110010 00010010 00001010 01101010 00111010 01011010 11001010 10101010 11111010 10000010 11100010
01000001 00100001 01110001 00010001 00001001 01101001 00111001 01011001 11001001 10101001 11111001 10110001 11010001
00000101 01100101 00110101 01010101 01001101 00101101 01111101 00011101 10001101 11011101 11001101 11110101 10010101
01000111 00100111 00110111 00010111 00001111 01101111 00111111 01011111 11001111 10101111 11111111 10000111 11100111
00000011 01100011 00110011 01010011 01001011 00101011 01111011 00011011 10001011 11101011 10111011 11100011 10010011
11001000 10101000 11111000 10011000 10000000 11100000 10110000 11010000 01000000 00110000 00010000 00110000 01011000
10001100 11101100 10111100 11011100 11000100 10100100 11110100 10010100 00000100 01100100 01010100 01011100 00011100
11001110 10101110 10011110 10000110 11100110 10110110 10110110 11010110 01000110 00001110 00001110 01111110 01011110
10001010 11101010 10110101 11000010 10100010 11110010 10010010 00000010 00000010 01100010 01010101 01110101 00011010
10001001 11101001 10111001 11000001 10100001 11110001 10110001 10010001 00000001 01100001 01010001 01111001 00011001
11001101 10101101 11111101 10011101 10000101 11100101 10110101 11010101 01000101 01000101 00010101 01110101 01011101
10001111 11101111 10111111 11001111 11000111 10100111 11110111 10010111 00000111 01100111 01010111 01111111 00011111
11001011 11111011 10011011 10011011 11100011 11100011 10110011 10110011 11010011 11010011 00100011 01110011 01011011

```

FIG. 11B

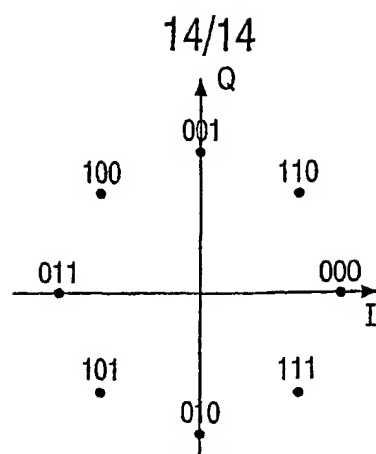


FIG. 12A

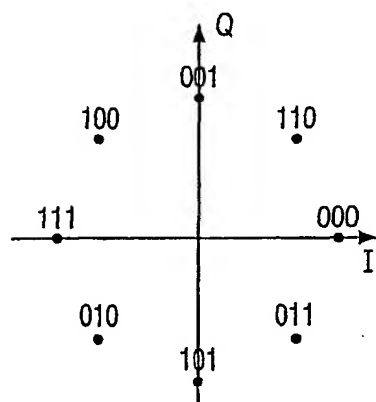


FIG. 12B

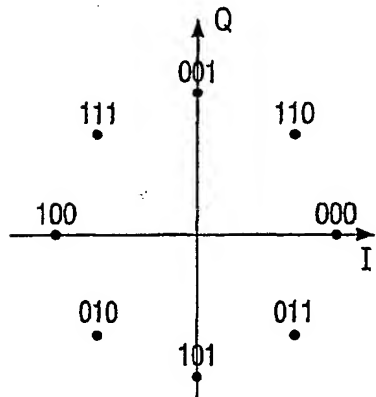


FIG. 12C

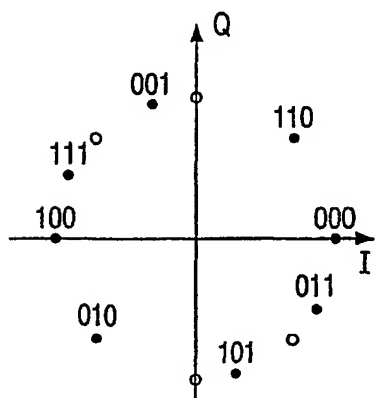


FIG. 13

## INTERNATIONAL SEARCH REPORT

Int. Application No

PCT/IB 01/02556

A. CLASSIFICATION OF SUBJECT MATTER  
IPC 7 H04L27/34 H03M13/00

According to International Patent Classification (IPC) or to both national classification and IPC

## B. FIELDS SEARCHED

Minimum documentation searched (classification system followed by classification symbols)  
IPC 7 H04L H03M

Documentation searched other than minimum documentation to the extent that such documents are included in the fields searched

Electronic data base consulted during the international search (name of data base and, where practical, search terms used)

EPO-Internal, PAJ, WPI Data

## C. DOCUMENTS CONSIDERED TO BE RELEVANT

Category *	Citation of document, with indication, where appropriate, of the relevant passages	Relevant to claim No.
A	LI X ET AL: "TRELLIS-CODED MODULATION WITH BIT INTERLEAVING AND ITERATIVE DOCODING" IEEE JOURNAL ON SELECTED AREAS IN COMMUNICATIONS, IEEE INC. NEW YORK, US, vol. 17, no. 4, April 1999 (1999-04), pages 715-724, XP000824314 ISSN: 0733-8716 cited in the application page 718, right-hand column, paragraph D	1,8,11, 14,17, 20,23
A	EP 0 948 140 A (LUCENT TECHNOLOGIES INC) 6 October 1999 (1999-10-06) cited in the application column 5, line 5 - line 7 ----- -/--	1,8,11, 14,17, 20,23

☒ Further documents are listed in the continuation of box C.

☒ Patent family members are listed in annex.

## \* Special categories of cited documents:

\*A\* document defining the general state of the art which is not considered to be of particular relevance

\*E\* earlier document but published on or after the international filing date

\*L\* document which may throw doubts on priority claim(s) or which is cited to establish the publication date of another citation or other special reason (as specified)

\*O\* document referring to an oral disclosure, use, exhibition or other means

\*P\* document published prior to the international filing date but later than the priority date claimed

\*T\* later document published after the international filing date or priority date and not in conflict with the application but cited to understand the principle or theory underlying the invention

\*X\* document of particular relevance; the claimed invention cannot be considered novel or cannot be considered to involve an inventive step when the document is taken alone

\*Y\* document of particular relevance; the claimed invention cannot be considered to involve an inventive step when the document is combined with one or more other such documents, such combination being obvious to a person skilled in the art.

\*Z\* document member of the same patent family

Date of the actual completion of the international search

26 March 2002

Date of mailing of the international search report

12/04/2002

Name and mailing address of the ISA

European Patent Office, P.B. 5818 Patentlaan 2  
NL - 2280 HV Rijswijk  
Tel. (+31-70) 340-2040, Tx. 31 651 epo nl,  
Fax: (+31-70) 340-3016

Authorized officer

Farese, L

## INTERNATIONAL SEARCH REPORT

Int. Patent Application No.

PCT/IB 01/02556

C.(Continuation) DOCUMENTS CONSIDERED TO BE RELEVANT		
Category *	Citation of document, with indication, where appropriate, of the relevant passages	Relevant to claim No.
P,X	GOROKHOV A ; VAN DIJK M: "Optimised labeling maps for bit-interleaved transmission with turbo demodulation" IEEE VTS 53RD VEHICULAR TECHNOLOGY CONFERENCE, PROCEEDINGS, 6 - 9 May 2001, pages 1459-1463, XP002194006 Rhodes, Greece the whole document -----	1-29

Form PCT/ISA/210 (continuation of second sheet) (July 1992)

## FURTHER INFORMATION CONTINUED FROM PCT/ISA/ 210

Continuation of Box I.2

Claims Nos.: 4-7, 26-29

Present claims 4-7, 26-29 relate to an extremely large number of possible constellations. Support within the meaning of Article 6 PCT and disclosure within the meaning of Article 5 PCT is to be found, however, for only a very small proportion of the constellations claimed. In the present case, the claims so lack support, and the application so lacks disclosure, that a meaningful search over the whole of the claimed scope is impossible. Consequently, the search has been carried out for those parts of the claims which appear to be supported and disclosed, namely those parts relating to the constellation corresponding to figures 8A-8G, 9A-9C, 11A-11C, 12A-12C.

The applicant's attention is drawn to the fact that claims, or parts of claims, relating to inventions in respect of which no international search report has been established need not be the subject of an international preliminary examination (Rule 66.1(e) PCT). The applicant is advised that the EPO policy when acting as an International Preliminary Examining Authority is normally not to carry out a preliminary examination on matter which has not been searched. This is the case irrespective of whether or not the claims are amended following receipt of the search report or during any Chapter II procedure.

## INTERNATIONAL SEARCH REPORT

information on patent family members

Initial Application No

PCT/IB 01/02556

Patent document cited in search report		Publication date	Patent family member(s)	Publication date
EP 0948140	A	06-10-1999	EP 0948140 A1	06-10-1999
			AU 721048 B2	22-06-2000
			AU 2251099 A	14-10-1999
			BR 9901231 A	18-01-2000
			CN 1236229 A	24-11-1999
			JP 2000041078 A	08-02-2000
			US 6353911 B1	05-03-2002

Form PCT/ISA/210 (patent family annex) (July 1992)

